

# A NOTE ON COMPLETE SUBDIVISIONS IN DIGRAPHS OF LARGE OUTDEGREE

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ABSTRACT. Mader conjectured that for all  $\ell$  there is an integer  $\delta^+(\ell)$  such that every digraph of minimum outdegree at least  $\delta^+(\ell)$  contains a subdivision of a transitive tournament of order  $\ell$ . In this note we observe that if the minimum outdegree of a digraph is sufficiently large compared to its order then one can even guarantee a subdivision of a large complete digraph. More precisely, let  $\vec{G}$  be a digraph of order  $n$  whose minimum outdegree is at least  $d$ . Then  $\vec{G}$  contains a subdivision of a complete digraph of order  $\lfloor d^2/(8n^{3/2}) \rfloor$ .

## 1. INTRODUCTION

A fundamental result of Mader [4] states that for every integer  $\ell$  there is a smallest  $d = d(\ell)$  so that every graph of average degree at least  $d$  contains a subdivision of a complete graph on  $\ell$  vertices. Bollobás and Thomason [1] as well as Komlós and Szemerédi [3] showed that  $d(\ell)$  is quadratic in  $\ell$ . In [6], Mader made the following conjecture, which would provide a digraph analogue of these results (a transitive tournament is a complete graph whose edges are oriented transitively).

**Conjecture 1 (Mader [6]).** *For every integer  $\ell > 0$  there is a smallest integer  $\delta^+(\ell)$  such that every digraph  $\vec{G}$  with minimum outdegree at least  $\delta^+(\ell)$  contains a subdivision of the transitive tournament on  $\ell$  vertices.*

It is easy to see that  $\delta^+(\ell) = \ell - 1$  for  $\ell \leq 3$ . Mader [7] showed that  $\delta^+(4) = 3$ . Even the existence of  $\delta^+(5)$  is not known. One might be tempted to conjecture that large minimum outdegree would even force the existence of a subdivision of a large complete digraph (a complete digraph has a directed edge from  $v$  to  $w$  for any ordered pair  $v, w$  of vertices). However, for all  $n$  Thomassen [9] constructed a digraph on  $n$  vertices whose minimum outdegree is at least  $\frac{1}{2} \log_2 n$  but which does not contain an even directed cycle (and thus no complete digraph on 3 vertices). The additional assumption of large minimum indegree in Conjecture 1 does not help either. Mader [6] modified the construction in [9] to obtain digraphs having arbitrarily large minimum indegree and outdegree without a subdivision of a complete digraph on 3 vertices.

The fact that one certainly cannot replace the minimum outdegree in Conjecture 1 by the average degree is easy to see: consider the complete bipartite graph with equal size vertex classes and orient all edges from the first to the second class. The resulting digraph  $\vec{B}$  has average degree  $|\vec{B}|/2$  but not even a directed cycle or a transitive tournament on 3 vertices. (On the other hand, Jagger [2] showed that if the average degree of a digraph  $\vec{G}$  is a little larger than  $|\vec{G}|/2$ , then  $\vec{G}$  does contain a subdivision of a large complete digraph.)

So in some sense, the above examples and constructions show that Conjecture 1 is the only possible analogue of the result in [4] mentioned above. Our main result is that if the minimum outdegree of a digraph is sufficiently large compared to its order, then Conjecture 1 is true. In fact, we show that in this case, one can even guarantee a subdivision of a complete digraph.

**Theorem 2.** *Let  $\vec{G}$  be a digraph of order  $n$  whose minimum outdegree is at least  $d$ . Then  $\vec{G}$  contains a subdivision of the complete digraph of order  $\lfloor d^2/(8n^{3/2}) \rfloor$ .*

Note that the bound is nontrivial as soon as  $d$  is a little larger than  $n^{3/4}$ . Also, recall that the result of Thomassen [9] mentioned above implies that we cannot have a subdivision of a complete digraph of order at least 3 if  $d \leq \frac{1}{2} \log_2 n$ . Furthermore, note that if  $d = cn$ , then Theorem 2 guarantees a subdivision of a complete digraph of order  $\lfloor c' \sqrt{n} \rfloor$ , where  $c' = c^2/8$ . It is easy to see that this is best possible up to the value of  $c'$  (consider the complete bipartite digraph with vertex classes of equal size).

The main ingredient in the proof of Theorem 2 is Lemma 4. It states that if  $\vec{G}$  has  $n$  vertices and its minimum outdegree is  $\gg \sqrt{n}$ , then  $\vec{G}$  has a subdigraph  $\vec{H}$  which is highly connected in the following sense: if  $x$  is any vertex of  $\vec{H}$  and  $y$  is a vertex of large indegree, then there are many internally disjoint dipaths from  $x$  to  $y$  in  $\vec{H}$ . Lemma 4 also guarantees the existence of many such vertices  $y$ . For undirected graphs, there is a much stronger result of Mader [5] which implies that every graph of minimum degree at least  $4k$  has a  $k$ -connected subgraph. Since a digraph version of this result is not known, Lemma 4 may be of independent interest. There are also several related results of Mader [6, 8] which investigate the existence of pairs of vertices with large local connectivity in digraphs of large minimum outdegree. The proof of Lemma 4 is quite elementary: if the current subdigraph  $\vec{H}$  does not satisfy the requirements, then we can use Menger's theorem to find a significantly smaller subdigraph whose minimum outdegree is almost as large as that of  $\vec{H}$ . Since this means that the density of the successive subdigraphs increases, this process must eventually terminate.

## 2. PROOF OF THEOREM 2

Before we start with the proof of Theorem 2 let us introduce some notation. The digraphs  $\vec{G}$  considered in this note do not contain loops and between any ordered vertex pair  $x, y \in \vec{G}$  there is at most one edge from  $x$  to  $y$ . (There might also be another edge from  $y$  to  $x$ .) We denote by  $\delta^+(\vec{G})$  the minimum outdegree of a digraph  $\vec{G}$  and by  $|\vec{G}|$  its order. We write  $d_{\vec{G}}^+(x)$  for the outdegree of a vertex  $x \in \vec{G}$  and  $d_{\vec{G}}^-(x)$  for its indegree. A digraph  $\vec{H}$  is a *subdivision* of  $\vec{G}$  if  $\vec{H}$  can be obtained from  $\vec{G}$  by replacing each edge  $\vec{xy} \in \vec{G}$  with a dipath from  $x$  to  $y$  such that all these dipaths are internally disjoint for distinct edges. The vertices of  $\vec{H}$  corresponding to the vertices of  $\vec{G}$  are called *branch vertices*.

Given two vertices  $x$  and  $y$  of a digraph  $\vec{G}$ , we define  $\kappa_{\vec{G}}(x, y)$  to be the largest integer  $1 \leq k \leq |\vec{G}| - 2$  such that  $\vec{G} - S$  contains a dipath from  $x$  to  $y$

for every vertex set  $S \subseteq V(\vec{G}) \setminus \{x, y\}$  of size  $< k$ . We define  $\kappa_{\vec{G}}(x, y) := 0$  if  $\vec{G}$  does not contain a dipath from  $x$  to  $y$ . We will use the following version of Menger's theorem for digraphs.

**Theorem 3 (Menger's theorem for digraphs).** *Let  $x$  and  $y$  be vertices of a digraph  $\vec{G}$  such that  $\kappa_{\vec{G}}(x, y) \geq k$ . Then  $\vec{G}$  contains  $k$  internally disjoint dipaths from  $x$  to  $y$ .*

As mentioned above, the main step in the proof of Theorem 2 is to find a subdigraph  $\vec{H}$  of  $\vec{G}$  such that the minimum outdegree of  $\vec{H}$  is still large and such that every vertex of  $\vec{H}$  sends many internally disjoint dipaths to each vertex of  $\vec{H}$  which has large indegree.

**Lemma 4.** *Every digraph  $\vec{G}$  of order  $n$  with  $\delta^+(\vec{G}) \geq d$  contains a subdigraph  $\vec{H}$  such that*

- (i)  $\delta^+(\vec{H}) > d/2$ ,
- (ii)  $\kappa_{\vec{H}}(x, y) \geq d^2/(4n)$  for all pairs  $x, y \in V(\vec{H})$  with  $d_{\vec{H}}^-(y) \geq d/2$ ,
- (iii) at least  $d^2/(4n)$  vertices of  $\vec{H}$  have indegree at least  $d/2$  in  $\vec{H}$ .

**Proof.** Put

$$\alpha := \frac{d}{n} \quad \text{and} \quad \alpha' := \frac{d^2}{4n^2} = \frac{\alpha^2}{4}.$$

By Theorem 3 we may assume that  $\kappa_{\vec{G}}(x, y) < \alpha'n$  for some vertices  $x, y$  of  $\vec{G}$  with  $d_{\vec{G}}^-(y) \geq d/2$ . Otherwise we could take  $\vec{H} := \vec{G}$ . (It is easy to check that  $\vec{H}$  then also satisfies condition (iii) of the lemma.) Let  $S \subseteq V(\vec{G}) \setminus \{x, y\}$  be a set of size  $< \alpha'n$  such that  $\vec{G} - S$  does not contain a dipath from  $x$  to  $y$ . Let  $Y$  be the set of all those vertices  $z$  for which  $\vec{G} - S$  contains a dipath from  $z$  to  $y$ . Then  $Y \cup S$  contains  $y$  as well as all the at least  $d/2 = \alpha n/2$  inneighbours of  $y$ . Let  $C$  denote the component of the undirected graph corresponding to  $\vec{G} - (Y \cup S)$  which contains  $x$ . Let  $\vec{G}_1$  be the subdigraph of  $\vec{G}$  induced by all vertices in  $C$ . Then  $|\vec{G}_1| \leq n - |Y \cup S| < (1 - \alpha/2)n$ . Moreover, note that there exists no edge directed from a vertex of  $\vec{G}_1$  to a vertex outside  $V(\vec{G}_1) \cup S$ . Thus

$$(1) \quad \delta^+(\vec{G}_1) \geq \delta^+(\vec{G}) - |S| > (\alpha - \alpha')n.$$

If  $\vec{G}_1$  does not satisfy condition (ii) of the lemma we again apply Theorem 3 to obtain a subdigraph  $\vec{G}_2 \subseteq \vec{G}_1$ . We continue in this fashion until we obtain a subdigraph  $\vec{G}_r$  which satisfies condition (ii). We will show that  $\vec{G}_r$  also satisfies (i) and (iii). Put  $\vec{G}_0 := \vec{G}$ ,

$$\delta_i := \frac{\delta^+(\vec{G}_i)}{|\vec{G}_i|} \quad \text{and} \quad \gamma_{i-1} := \frac{|\vec{G}_{i-1}|}{|\vec{G}_i|}$$

for all  $i \leq r$ . Similarly as in (1) it follows that

$$(2) \quad \delta^+(\vec{G}_i) = \delta_i |\vec{G}_i| \geq \delta_{i-1} |\vec{G}_{i-1}| - \alpha'n \geq (\alpha - i\alpha')n.$$

Thus  $\delta_i \geq \delta_{i-1}\gamma_{i-1} - \alpha'n/|\vec{G}_i| = \delta_{i-1}\gamma_{i-1} - \alpha' \prod_{j=0}^{i-1} \gamma_j$ . Using this inequality and induction on  $i$  one can show that

$$(3) \quad \delta_i \geq (\alpha - i\alpha') \prod_{j=0}^{i-1} \gamma_j = (\alpha - i\alpha') \frac{n}{|\vec{G}_i|}.$$

Since we delete at least  $d/2 = \alpha n/2$  vertices when going from  $\vec{G}_{i-1}$  to  $\vec{G}_i$  (namely the inneighbours of the vertex playing the role of  $y$ ), we have that  $|\vec{G}_r| \leq n - r\alpha n/2$ . In particular this shows that  $r < 2/\alpha$ . However, since (3) implies that  $1 > \delta_r \geq (\alpha - r\alpha')/(1 - r\alpha/2)$  we even have  $r < (1 - \alpha)/(\alpha/2 - \alpha')$ . Thus

$$(4) \quad \delta^+(\vec{G}_i) \stackrel{(2)}{\geq} (\alpha - r\alpha')n \geq \left(\alpha - \frac{1 - \alpha}{2/\alpha - 1}\right)n = \frac{\alpha n}{2 - \alpha} > \frac{d}{2}.$$

Altogether this shows that  $\vec{G}_r =: \vec{H}$  satisfies conditions (i) and (ii) of the lemma. To check that  $\vec{H}$  also satisfies condition (iii) let  $\ell$  denote the number of vertices of indegree  $\geq d/2$  in  $\vec{H}$ . Then

$$\frac{\alpha n |\vec{H}|}{2 - \alpha} \stackrel{(4)}{\leq} \delta^+(\vec{H}) |\vec{H}| \leq |\vec{H}| \frac{d}{2} + \ell |\vec{H}|,$$

which implies that  $\ell \geq \alpha d/(4 - 2\alpha) \geq d^2/(4n)$ , as required.  $\square$

**Proof of Theorem 2.** Let  $\ell := \lfloor d^2/(8n^{3/2}) \rfloor$ . We first apply Lemma 4 to obtain a subdigraph  $\vec{H} \subseteq \vec{G}$  as described there. We pick a set  $X \subseteq V(\vec{H})$  of  $\ell$  vertices having indegree  $\geq d/2$  in  $\vec{H}$ . (Such a set  $X$  exists by condition (iii) of Lemma 4.)  $X$  will be the set of our branch vertices. For every pair  $x, y \in X$  there exist at least  $d^2/(4n)$  internally disjoint dipaths from  $x$  to  $y$ . Thus the average number of inner vertices on such a path is at most  $4n^2/d^2$ . Hence  $\vec{H}$  contains at least  $d^2/(8n)$  internally disjoint dipaths from  $x$  to  $y$  such that each of these has at most  $8n^2/d^2$  inner vertices. Let us call such a dipath *short*. This shows that we can connect all pairs of branch vertices greedily (in both directions) by choosing each time a short dipath which is internally disjoint from all the short dipaths chosen before. In each step we destroy at most  $8n^2/d^2$  further dipaths. But  $(|X|^2 - 1)8n^2/d^2 < 8\ell^2 n^2/d^2 \leq d^2/(8n)$ , so we can connect all pairs of branch vertices by short dipaths.  $\square$

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